

CS 114
Introduction to Computational Linguistics

Grammar and Parsing (II)

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Thanks to Dan Jurafsky and Jim Martin for many of these slides!

www.cs.brandeis.edu/~cs114/slides/114.08.lec10.ppt

www.stanford.edu/class/linguist180/180.07.lec10.ppt

Grammars and Parsing

- Context-Free Grammars and Constituency
- Some common CFG phenomena for English
- Baby parsers: Top-down and Bottom-up Parsing
- Today: Real parsers: Dynamic Programming parsing
 - CKY
- Probabilistic parsing
- Optional section: the Earley algorithm

Dynamic Programming

- We need a method that fills a table with partial results that
 - Does not do (avoidable) repeated work
 - Does not fall prey to left-recursion
 - Can find all the pieces of an exponential number of trees in polynomial time.
- Two popular methods
 - CKY
 - Earley

The CKY (Cocke-Kasami-Younger) Algorithm

- Requires the grammar be in Chomsky Normal Form (CNF)
 - All rules must be in following form:
 - $A \rightarrow BC$
 - $A \rightarrow w$
- Any grammar can be converted automatically to Chomsky Normal Form

Converting to CNF

- Rules that mix terminals and non-terminals
 - Introduce a new dummy non-terminal that covers the terminal
 - INFVP \rightarrow to VP replaced by:
 - INFVP \rightarrow TO VP
 - TO \rightarrow to
- Rules that have a single non-terminal on right ("unit productions")
 - Rewrite each unit production with the RHS of their expansions
- Rules whose right hand side length >2
 - Introduce dummy non-terminals that spread the right-hand side

Automatic Conversion to CNF

$S \rightarrow NP VP$	$S \rightarrow NP VP$
$S \rightarrow Aux NP VP$	$S \rightarrow XI VP$
	$XI \rightarrow Aux NP$
$S \rightarrow VP$	$S \rightarrow book \mid include \mid prefer$
	$S \rightarrow Verb NP$
	$S \rightarrow VP PP$
$NP \rightarrow Det Nominal$	$NP \rightarrow Det Nominal$
$NP \rightarrow Proper-Noun$	$NP \rightarrow TWA \mid Houston$
$NP \rightarrow Pronoun$	$NP \rightarrow I \mid she \mid me$
$Nominal \rightarrow Noun$	$Nominal \rightarrow book \mid flight \mid meal \mid money$
$Nominal \rightarrow Noun Nominal$	$Nominal \rightarrow Noun Nominal$
$Nominal \rightarrow Nominal PP$	$Nominal \rightarrow Nominal PP$
$VP \rightarrow Verb$	$VP \rightarrow book \mid include \mid prefer$
$VP \rightarrow Verb NP$	$VP \rightarrow Verb NP$
$VP \rightarrow VP PP$	$VP \rightarrow VP PP$
$PP \rightarrow Prep NP$	$PP \rightarrow Prep NP$

Figure 10.15 Original L0 Grammar and its conversion to CNF

Sample Grammar

<i>S</i> → <i>NP VP</i>	<i>Det</i> → <i>that</i> <i>this</i> <i>a</i>
<i>S</i> → <i>Aux NP VP</i>	<i>Noun</i> → <i>book</i> <i>flight</i> <i>meal</i> <i>money</i>
<i>S</i> → <i>VP</i>	<i>Verb</i> → <i>book</i> <i>include</i> <i>prefer</i>
<i>NP</i> → <i>Pronoun</i>	<i>Pronoun</i> → <i>I</i> <i>she</i> <i>me</i>
<i>NP</i> → <i>Proper-Noun</i>	<i>Proper-Noun</i> → <i>Houston</i> <i>TWA</i>
<i>NP</i> → <i>Det Nominal</i>	<i>Aux</i> → <i>does</i>
<i>Nominal</i> → <i>Noun</i>	<i>Preposition</i> → <i>from</i> <i>to</i> <i>on</i> <i>near</i> <i>through</i>
<i>Nominal</i> → <i>Nominal Noun</i>	
<i>Nominal</i> → <i>Nominal PP</i>	
<i>VP</i> → <i>Verb</i>	
<i>VP</i> → <i>Verb NP</i>	
<i>VP</i> → <i>Verb NP PP</i>	
<i>VP</i> → <i>Verb PP</i>	
<i>VP</i> → <i>VP PP</i>	
<i>PP</i> → <i>Preposition NP</i>	

Back to CKY Parsing

- Given rules in CNF
- Consider the rule $A \rightarrow BC$
 - If there is an A in the input then there must be a B followed by a C in the input.
 - If the A goes from i to j in the input then there must be some k st. $i < k < j$
 - Ie. The B splits from the C someplace.

CKY

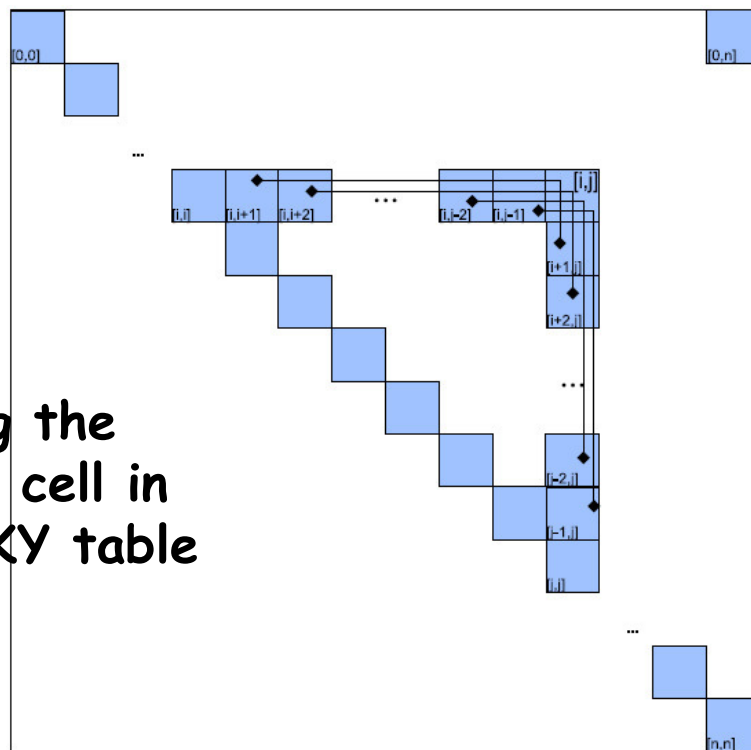
- So let's build a table so that an A spanning from i to j in the input is placed in cell $[i, j]$ in the table.
- So a non-terminal spanning an entire string will sit in cell $[0, n]$
- If we build the table bottom up we'll know that the parts of the A must go from i to k and from k to j

CKY

- Meaning that for a rule like $A \rightarrow B C$ we should look for a B in $[i,k]$ and a C in $[k,j]$.
- In other words, if we think there might be an A spanning i,j in the input... AND
- $A \rightarrow B C$ is a rule in the grammar THEN
- There must be a B in $[i,k]$ and a C in $[k,j]$ for some $i < k < j$
- So just loop over the possible k values

CKY Table

• Filling the $[i,j]$ th cell in the CKY table



CKY Algorithm

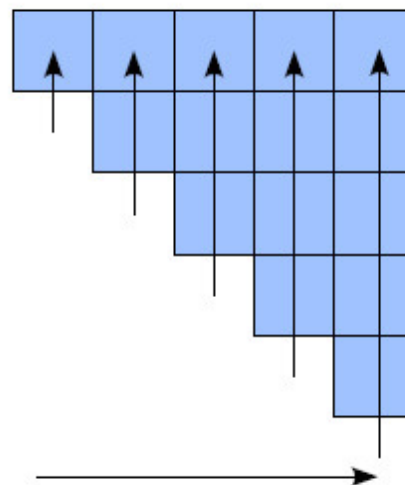
```
function CKY-PARSE(words, grammar) returns table  
  
  for j ← from 1 to LENGTH(words) do  
    table[j - 1, j] ← {A | A → words[j] ∈ grammar }  
    for i ← from j - 2 downto 0 do  
      for k ← i + 1 to j - 1 do  
        table[i, j] ← table[i, j] ∪  
          {A | A → BC ∈ grammar,  
            B ∈ table[i, k],  
            C ∈ table[k, j] }
```

Note

- We arranged the loops to fill the table a column at a time, from left to right, bottom to top.
 - This assures us that whenever we're filling a cell, the parts needed to fill it are already in the table (to the left and below)
 - Are there other ways to fill the table?

0 Book 1 the 2 flight 3 through 4 Houston 5

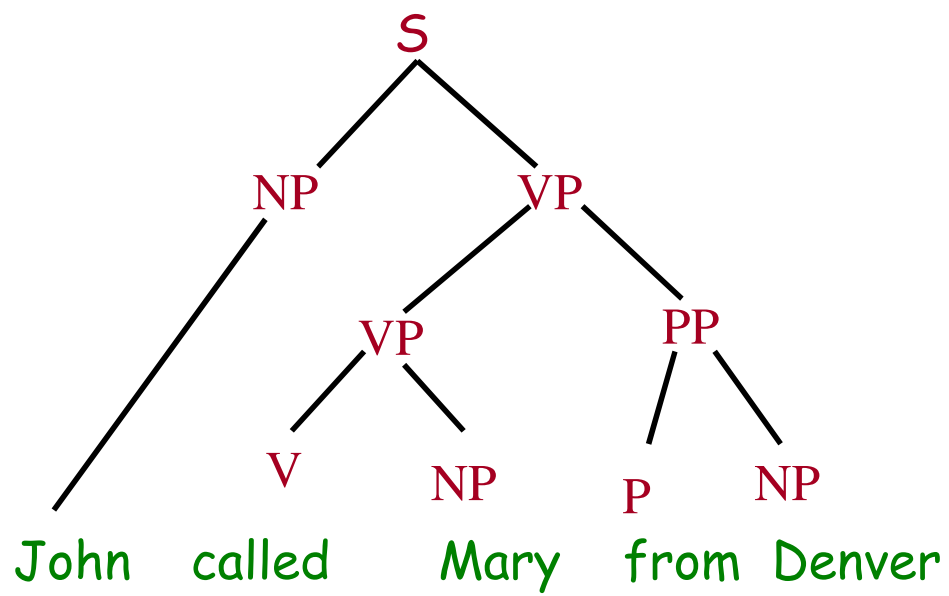
<i>Book</i>	<i>the</i>	<i>flight</i>	<i>through</i>	<i>Houston</i>
S,VP,Verb Nominal, Noun [0,1]		S,VP,X2 [0,3]		S, VP [0,5]
	Det [1,2]	NP [1,3]		NP [1,5]
		Nominal, Noun [2,3]		Nominal [2,5]
			Prep [3,4]	PP [3,5]
				NP, Proper- Noun [4,5]



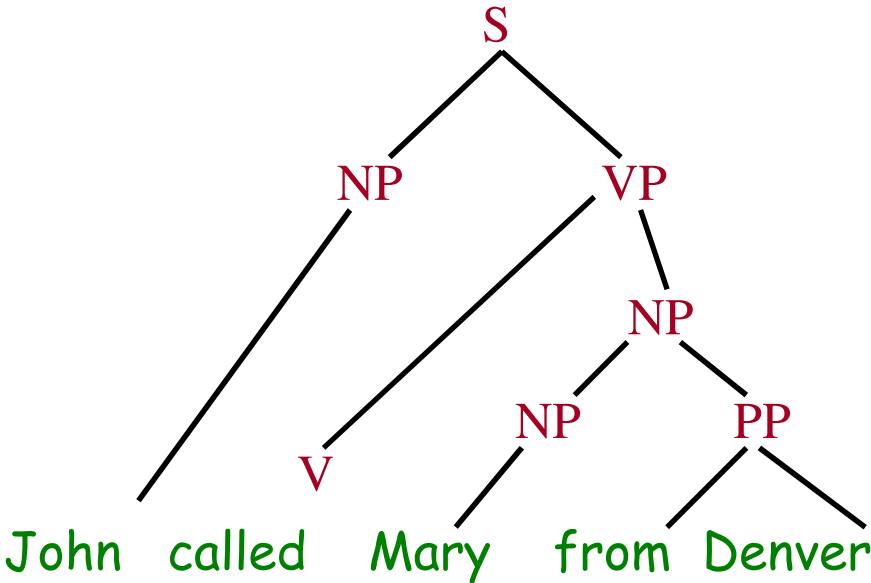
CYK Example

- S -> NP VP
- VP -> V NP
- NP -> NP PP
- VP -> VP PP
- PP -> P NP
- NP -> John, Mary, Denver
- V -> called
- P -> from

Example



Example



Example

				NP
			P	Denver
		NP	from	
	V	Mary		
NP	called			
John				

Example

				NP
			P	Denver
		NP	from	
X	V	Mary		
NP	called			
John				

Example

				NP
			P	Denver
	VP →	NP	from	
X	↓ V	Mary		
NP	called			
John				

Example

				NP
		X	P	Denver
	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

			PP	NP
		X	P	Denver
	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

			PP	NP
		X	P	Denver
S →	VP	NP	from	
↓	V	Mary		
NP	called			
John				

Example

			PP	NP
	X	X	P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

		NP	→ PP	NP
	X		P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

		NP	PP	NP
X	X	X	P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

	VP	NP	PP	NP
X	X	X	P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

	VP	NP	PP	NP
X	X	X	P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

	VP ₁ VP ₂	NP	PP	NP
X	X	X	P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

Example

S	VP ₁ VP ₂	NP	PP	NP
X	X	X	P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

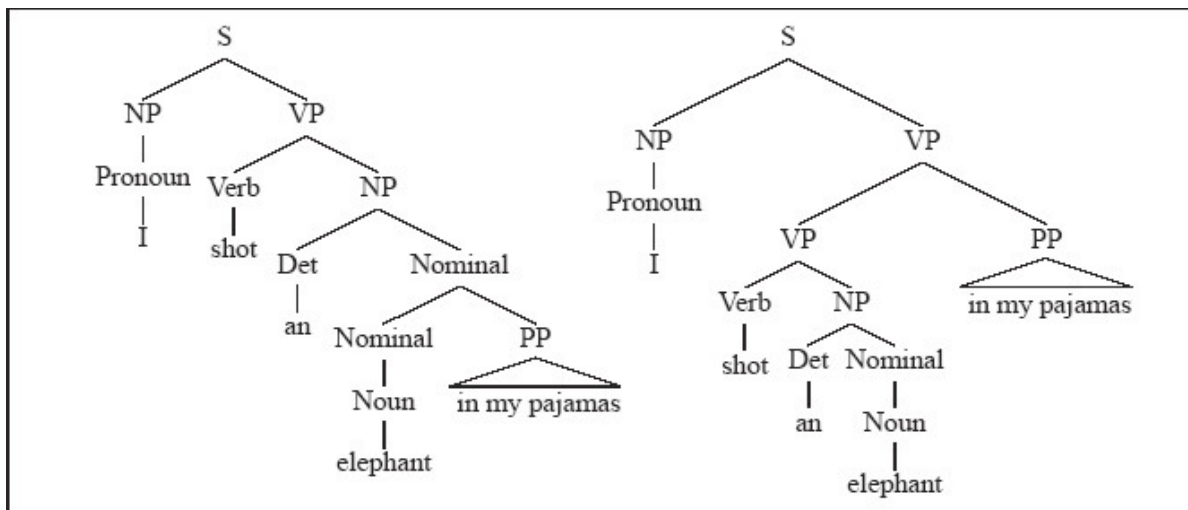
Example

S	VP	NP	PP	NP
X	X	X	P	Denver
S	VP	NP	from	
X	V	Mary		
NP	called			
John				

Back to Ambiguity

- Did we solve it?

Ambiguity



Ambiguity

- No...
 - Both CKY and Earley will result in multiple **S** structures for the $[0,n]$ table entry.
 - They both efficiently store the sub-parts that are shared between multiple parses.
 - But neither can tell us which one is right.
 - Not a parser – a recognizer
 - The presence of an S state with the right attributes in the right place indicates a successful recognition.
 - But no parse tree... no parser
 - That's how we solve (not) an exponential problem in polynomial time

Converting CKY from Recognizer to Parser

- With the addition of a few pointers we have a parser
- Augment each new cell in chart to point to where we came from.

Optional section: the Earley alg

Problem (minor)

- We said CKY requires the grammar to be binary (ie. In Chomsky-Normal Form).
- We showed that any arbitrary CFG can be converted to Chomsky-Normal Form so that's not a huge deal
- **Except** when you change the grammar the trees come out wrong
- All things being equal we'd prefer to leave the grammar alone.

Earley Parsing

- Allows arbitrary CFGs
- Where CKY is bottom-up, Earley is top-down
- Fills a table in a single sweep over the input words
 - Table is length $N+1$; N is number of words
 - Table entries represent
 - Completed constituents and their locations
 - In-progress constituents
 - Predicted constituents

States

- The table-entries are called states and are represented with dotted-rules.

S -> ' VP

A VP is predicted

NP -> Det ' Nominal

An NP is in progress

VP -> V NP '

A VP has been found

States/Locations

- It would be nice to know where these things are in the input
SO...

S -> ' VP [0,0]

A VP is predicted at the
start of the sentence

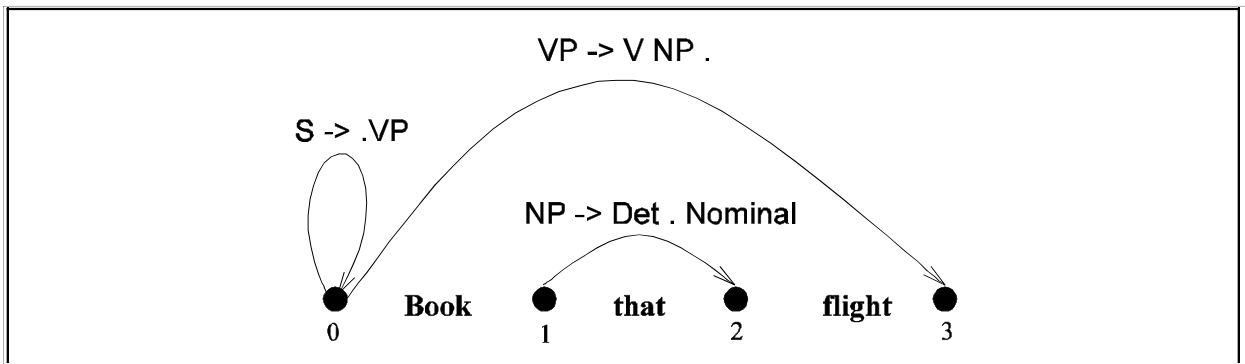
NP -> Det ' Nominal [1,2]

An NP is in progress; the
Det goes from 1 to 2

VP -> V NP ' [0,3]

A VP has been found
starting at 0 and ending at 3

Graphically



Earley

- As with most dynamic programming approaches, the answer is found by looking in the table in the right place.
- In this case, there should be an S state in the final column that spans from 0 to $n+1$ and is complete.
- If that's the case you're done.
 - $S \rightarrow \alpha \cdot [0, n+1]$

Earley Algorithm

- March through chart left-to-right.
- At each step, apply 1 of 3 operators
 - Predictor
 - Create new states representing top-down expectations
 - Scanner
 - Match word predictions (rule with word after dot) to words
 - Completer
 - When a state is complete, see what rules were looking for that completed constituent

Predictor

- Given a state
 - With a non-terminal to right of dot
 - That is not a part-of-speech category
 - Create a new state for each expansion of the non-terminal
 - Place these new states into same chart entry as generated state, beginning and ending where generating state ends.
 - So predictor looking at
 - S -> . VP [0,0]
 - results in
 - VP -> . Verb [0,0]
 - VP -> . Verb NP [0,0]

Scanner

- Given a state
 - With a non-terminal to right of dot
 - That is a part-of-speech category
 - If the next word in the input matches this part-of-speech
 - Create a new state with dot moved over the non-terminal
 - So scanner looking at
 - VP -> . Verb NP [0,0]
 - If the next word, "book", can be a verb, add new state:
 - VP -> Verb . NP [0,1]
 - Add this state to chart entry following current one
 - Note: Earley algorithm uses top-down input to disambiguate POS!
Only POS predicted by some state can get added to chart!

Completer

- Applied to a state when its dot has reached right end of rule.
- Parser has discovered a category over some span of input.
- Find and advance all previous states that were looking for this category
 - copy state, move dot, insert in current chart entry
- Given:
 - NP -> Det Nominal . [1,3]
 - VP -> Verb. NP [0,1]
- Add
 - VP -> Verb NP . [0,3]

Earley: how do we know we are done?

- How do we know when we are done?.
- Find an S state in the final column that spans from 0 to $n+1$ and is complete.
- If that's the case you're done.
 - $S \rightarrow \alpha \cdot [0, n+1]$

Earley

- So sweep through the table from 0 to $n+1$...
 - New predicted states are created by starting top-down from S
 - New incomplete states are created by advancing existing states as new constituents are discovered
 - New complete states are created in the same way.

Earley

- More specifically...
 1. Predict all the states you can upfront
 2. Read a word
 1. Extend states based on matches
 2. Add new predictions
 3. Go to 2
 3. Look at $N+1$ to see if you have a winner

Example

- Book that flight
- We should find... an S from 0 to 3 that is a completed state...

Example

Chart[0]	S0	$\gamma \rightarrow \bullet S$	[0,0]	Dummy start state
	S1	$S \rightarrow \bullet NP VP$	[0,0]	Predictor
	S2	$S \rightarrow \bullet Aux NP VP$	[0,0]	Predictor
	S3	$S \rightarrow \bullet VP$	[0,0]	Predictor
	S4	$NP \rightarrow \bullet Pronoun$	[0,0]	Predictor
	S5	$NP \rightarrow \bullet Proper-Noun$	[0,0]	Predictor
	S6	$NP \rightarrow \bullet Det Nominal$	[0,0]	Predictor
	S7	$VP \rightarrow \bullet Verb$	[0,0]	Predictor
	S8	$VP \rightarrow \bullet Verb NP$	[0,0]	Predictor
	S9	$VP \rightarrow \bullet Verb NP PP$	[0,0]	Predictor
	S10	$VP \rightarrow \bullet Verb PP$	[0,0]	Predictor
	S11	$VP \rightarrow \bullet VP PP$	[0,0]	Predictor

Example

Chart[1]			
S12	<i>Verb</i> → <i>book</i> •	[0,1]	Scanner
S13	<i>VP</i> → <i>Verb</i> •	[0,1]	Completer
S14	<i>VP</i> → <i>Verb</i> • <i>NP</i>	[0,1]	Completer
S15	<i>VP</i> → <i>Verb</i> • <i>NP PP</i>	[0,0]	Predictor
S16	<i>VP</i> → <i>Verb</i> • <i>PP</i>	[0,0]	Predictor
S17	<i>S</i> → <i>VP</i> •	[0,1]	Completer
S18	<i>VP</i> → <i>VP</i> • <i>PP</i>	[0,1]	Completer
S19	<i>NP</i> → • <i>Pronoun</i>	[1,1]	Predictor
S20	<i>NP</i> → • <i>Proper-Noun</i>	[1,1]	Predictor
S21	<i>NP</i> → • <i>Det Nominal</i>	[1,1]	Predictor
S22	<i>PP</i> → • <i>Prep NP</i>	[1,1]	Predictor

Example

Chart[2]	S23	<i>Det</i> → <i>that</i> •	[1,2]	Scanner
	S24	<i>NP</i> → <i>Det</i> • <i>Nominal</i>	[1,2]	Completer
	S25	<i>Nominal</i> → • <i>Noun</i>	[2,2]	Predictor
	S26	<i>Nominal</i> → • <i>Nominal Noun</i>	[2,2]	Predictor
	S27	<i>Nominal</i> → • <i>Nominal PP</i>	[2,2]	Predictor
Chart[3]	S28	<i>Noun</i> → <i>flight</i> •	[2,3]	Scanner
	S29	<i>Nominal</i> → <i>Noun</i> •	[2,3]	Completer
	S30	<i>NP</i> → <i>Det Nominal</i> •	[1,3]	Completer
	S31	<i>Nominal</i> → <i>Nominal</i> • <i>Noun</i>	[2,3]	Completer
	S32	<i>Nominal</i> → <i>Nominal</i> • <i>PP</i>	[2,3]	Completer
	S33	<i>VP</i> → <i>Verb NP</i> •	[0,3]	Completer
	S34	<i>VP</i> → <i>Verb NP</i> • <i>PP</i>	[0,3]	Completer
	S35	<i>PP</i> → • <i>Prep NP</i>	[3,3]	Predictor
	S36	<i>S</i> → <i>VP</i> •	[0,3]	Completer

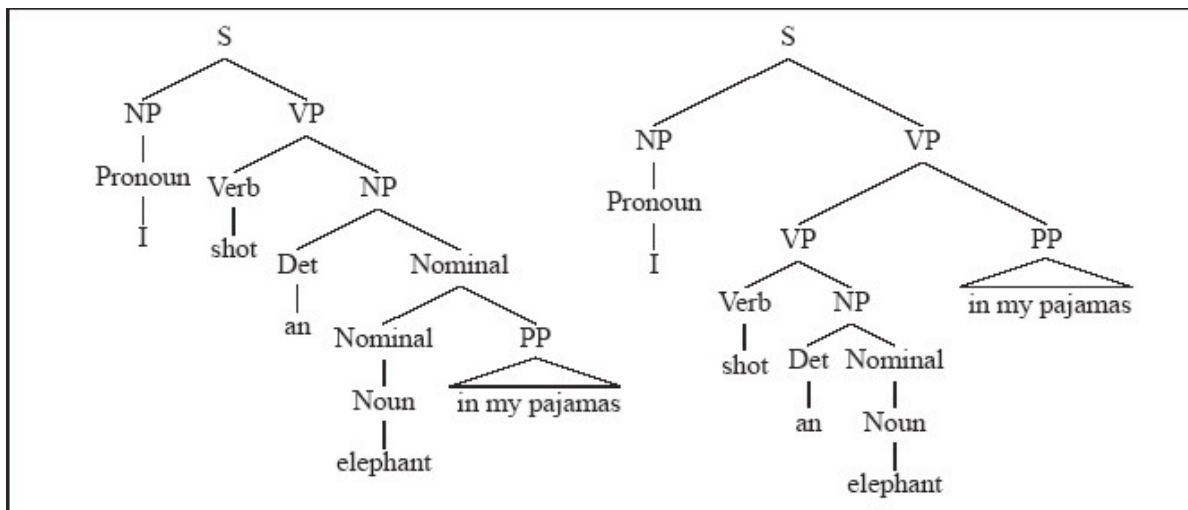
Details

- What kind of algorithms did we just describe (both Earley and CKY)
 - Not parsers – recognizers
 - The presence of an S state with the right attributes in the right place indicates a successful recognition.
 - But no parse tree... no parser
 - That's how we solve (not) an exponential problem in polynomial time

Back to Ambiguity

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Converting Earley from Recognizer to Parser

- With the addition of a few pointers we have a parser
- Augment the “Completer” to point to where we came from.

Augmenting the chart with structural information

Chart[1]				
S8	<i>Verb</i> ■ <i>book</i> ■	[0,1]	Scanner	
S9	<i>VP</i> ■ <i>Verb</i> ■	[0,1]	Completer	S8
S10	<i>S</i> ■ <i>VP</i> ■	[0,1]	Completer	S9
S11	<i>VP</i> ■ <i>Verb</i> ■ <i>NP</i>	[0,1]	Completer	S8
S12	<i>NP</i> ■ ■ <i>Det</i> <i>NOMINAL</i>	[1,1]	Predictor	
S13	<i>NP</i> ■ ■ <i>Proper-Noun</i>	[1,1]	Predictor	

Chart[2]			
<i>Det</i> ■ <i>that</i> ■	[1,2]	Scanner	
<i>NP</i> ■ <i>Det</i> ■ <i>NOMINAL</i>	[1,2]	Completer	
<i>NOMINAL</i> ■ ■ <i>Noun</i>	[2,2]	Predictor	
<i>NOMINAL</i> ■ ■ <i>Noun</i> <i>NOMINAL</i>	[2,2]	Predictor	

Retrieving Parse Trees from Chart

- All the possible parses for an input are in the table
- We just need to read off all the backpointers from every complete S in the last column of the table
- Find all the $S \rightarrow X \cdot [0, N+1]$
- Follow the structural traces from the Completer
- Of course, this won't be polynomial time, since there could be an exponential number of trees
- So we can at least represent ambiguity efficiently

How to do parse disambiguation

- Probabilistic methods
- Augment the grammar with probabilities
- Then modify the parser to keep only most probable parses
- And at the end, return the most probable parse

Probabilistic CFGs

- The probabilistic model
 - Assigning probabilities to parse trees
- Getting the probabilities for the model
- Parsing with probabilities
 - Slight modification to dynamic programming approach
 - Task is to find the max probability tree for an input

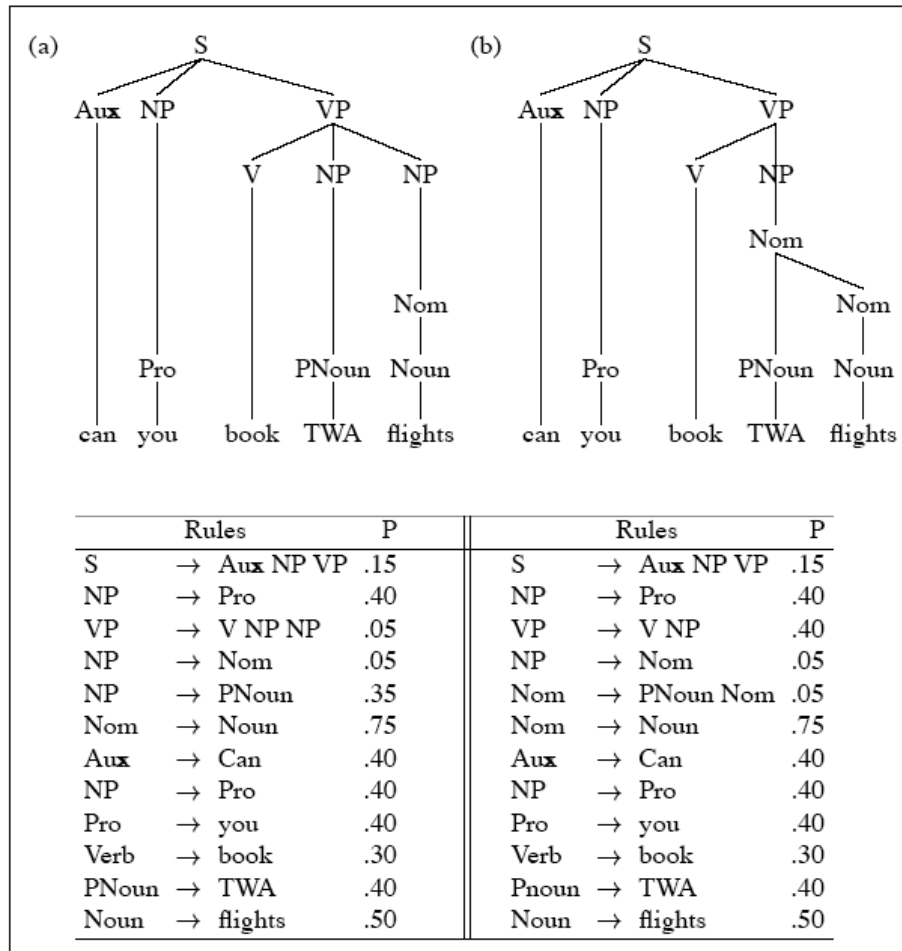
Probability Model

- Attach probabilities to grammar rules
 - The expansions for a given non-terminal sum to 1
- | | |
|------------------|-----|
| VP -> Verb | .55 |
| VP -> Verb NP | .40 |
| VP -> Verb NP NP | .05 |
- Read this as $P(\text{Specific rule} \mid \text{LHS})$

PCFG

$S \rightarrow NP VP$	[.80]	$Det \rightarrow that [.05] \mid the [.80] \mid a [.15]$
$S \rightarrow , Aux NP VP$	[.15]	$Noun \rightarrow , book$ [.10]
$S \rightarrow VP$	[.05]	$Noun \rightarrow flights$ [.50]
$NP \rightarrow Det Nom$	[.20]	$Noun \rightarrow meal$ [.40]
$NP \rightarrow Proper-Noun$	[.35]	$Verb \rightarrow book$ [.30]
$NP \rightarrow Nom$	[.05]	$Verb \rightarrow include$ [.30]
$NP \rightarrow Pronoun$	[.40]	$Verb \rightarrow want$ [.40]
$Nom \rightarrow Noun$	[.75]	$Aux \rightarrow can$ [.40]
$Nom \rightarrow Noun Nom$	[.20]	$Aux \rightarrow does$ [.30]
$Nom \rightarrow Proper-Noun Nom$	[.05]	$Aux \rightarrow do$ [.30]
$VP \rightarrow Verb$	[.55]	$Proper-Noun \rightarrow TWA$ [.40]
$VP \rightarrow Verb NP$	[.40]	$Proper-Noun \rightarrow Denver$ [.40]
$VP \rightarrow Verb NP NP$	[.05]	$Pronoun \rightarrow you [.40] \mid I [.60]$

PCFG



Probability Model (1)

- A derivation (tree) consists of the set of grammar rules that are in the tree
- The probability of a tree is just the product of the probabilities of the rules in the derivation.

Probability model

$$P(T, S) = \prod_{n \in T} p(r_n)$$

- $P(T, S) = P(T)P(S|T) = P(T)$; since $P(S|T)=1$

$$\begin{aligned} P(T_l) &= .15 * .40 * .05 * .05 * .35 * .75 * .40 * .40 * .40 \\ &\quad * .30 * .40 * .50 \\ &= 1.5 \times 10^{-6} \end{aligned}$$

$$\begin{aligned} P(T_r) &= .15 * .40 * .40 * .05 * .05 * .75 * .40 * .40 * .40 \\ &\quad * .30 * .40 * .50 \\ &= 1.7 \times 10^{-6} \end{aligned}$$

Probability Model (1.1)

- The probability of a word sequence $P(S)$ is the probability of its tree in the unambiguous case.
- It's the sum of the probabilities of the trees in the ambiguous case.

Getting the Probabilities

- From an annotated database (a treebank)
 - So for example, to get the probability for a particular VP rule just count all the times the rule is used and divide by the number of VPs overall.

TreeBanks

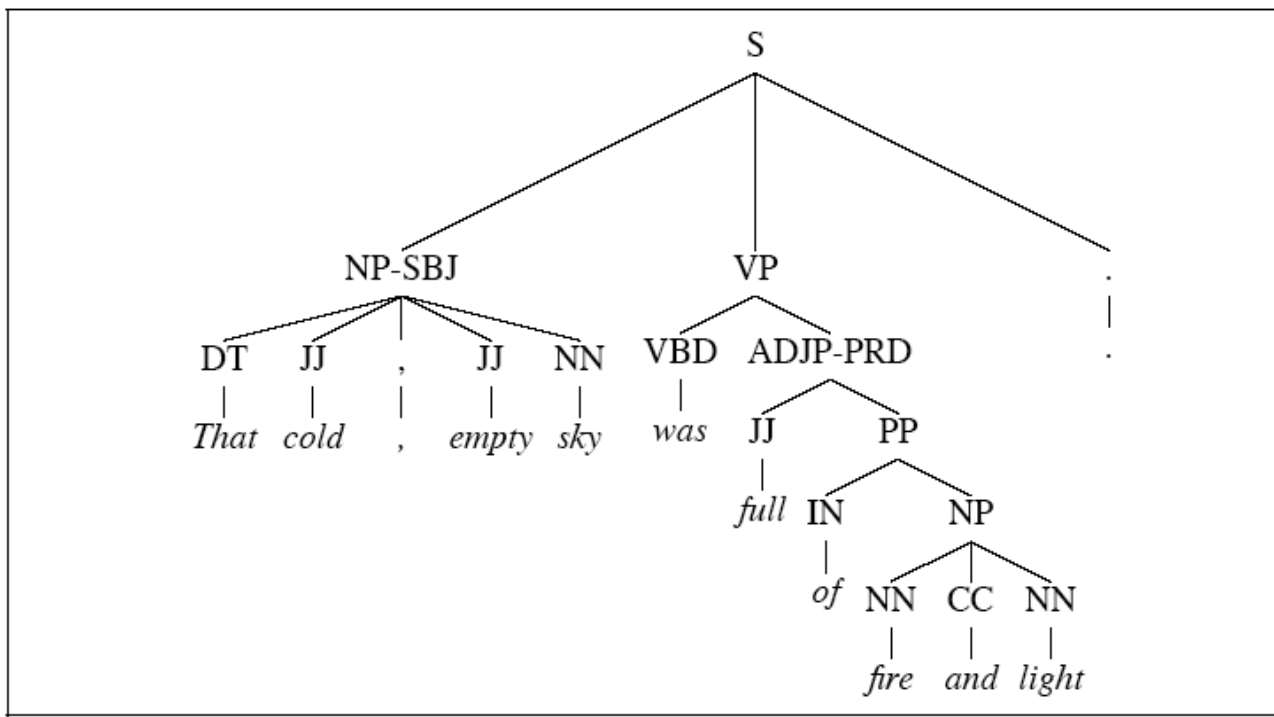
```
((S
  (NP-SBJ (DT That)
    (JJ cold) (, ,)
    (JJ empty) (NN sky) )
  (VP (VBD was)
    (ADJP-PRD (JJ full)
      (PP (IN of)
        (NP (NN fire)
          (CC and)
          (NN light) ))))
  (. .) ))
```

(a)

```
((S
  (NP-SBJ The/DT flight/NN )
  (VP should/MD
    (VP arrive/VB
      (PP-TMP at/IN
        (NP eleven/CD a.m/RB ))
      (NP-TMP tomorrow/NN ))))
```

(b)

Treebanks



Treebanks

```
( (S ( ' ' ' ' )
  (S-TPC-2
    (NP-SBJ-1 (PRP We) )
    (VP (MD would)
      (VP (VB have)
        (S
          (NP-SBJ (-NONE- *-1) )
          (VP (TO to)
            (VP (VB wait)
              (SBAR-TMP (IN until)
                (S
                  (NP-SBJ (PRP we) )
                  (VP (VBP have)
                    (VP (VBN collected)
                      (PP-CLR (IN on)
                        (NP (DT those) (NNS assets) ))))))))))))
    ( , , ) ( ' ' ' ' )
    (NP-SBJ (PRP he) )
    (VP (VBD said)
      (S (-NONE- *T*-2) ))
    ( . . ) ) )
```


Treebank Grammars

<i>S</i>	→ <i>NP VP .</i> <i>NP VP</i> " <i>S</i> ", <i>NP VP</i> . - <i>NONE</i> - <i>DT NN</i> <i>DT NN NNS</i> <i>NN CC NN</i> <i>CD RB</i>	<i>PRP</i>	→ <i>we he</i>
		<i>DT</i>	→ <i>the that those</i>
		<i>JJ</i>	→ <i>cold empty full</i>
		<i>NN</i>	→ <i>sky fire light flight</i>
		<i>NNS</i>	→ <i>assets</i>
		<i>CC</i>	→ <i>and</i>
		<i>IN</i>	→ <i>of at until on</i>
		<i>CD</i>	→ <i>eleven</i>
<i>NP</i>	→ <i>DT JJ , JJ NN</i> <i>PRP</i> - <i>NONE</i> -	<i>RB</i>	→ <i>a.m</i>
		<i>VB</i>	→ <i>arrive have wait</i>
<i>VP</i>	→ <i>MD VP</i> <i>VBD ADJP</i> <i>VBD S</i> <i>VB PP</i> <i>VB S</i> <i>VB SBAR</i> <i>VBP VP</i> <i>VBN VP</i> <i>TO VP</i>	<i>VBD</i>	→ <i>said</i>
		<i>VBP</i>	→ <i>have</i>
		<i>VBN</i>	→ <i>collected</i>
		<i>MD</i>	→ <i>should would</i>
		<i>TO</i>	→ <i>to</i>
<i>SBAR</i>	→ <i>IN S</i>		
<i>ADJP</i>	→ <i>JJ PP</i>		
<i>PP</i>	→ <i>IN NP</i>		

Lots of flat rules

NP → DT JJ NN
NP → DT JJ NNS
NP → DT JJ NN NN
NP → DT JJ JJ NN
NP → DT JJ CD NNS
NP → RB DT JJ NN NN
NP → RB DT JJ JJ NNS
NP → DT JJ JJ NNP NNS
NP → DT NNP NNP NNP NNP JJ NN
NP → DT JJ NNP CC JJ JJ NN NNS
NP → RB DT JJS NN NN SBAR
NP → DT VBG JJ NNP NNP CC NNP
NP → DT JJ NNS , NNS CC NN NNS NN
NP → DT JJ JJ VBG NN NNP NNP FW NNP
NP → NP JJ , JJ ' ' SBAR ' ' NNS

Example sentences from those rules

- Total: over 17,000 different grammar rules in the 1-million word Treebank corpus

(9.19) [DT The] [JJ state-owned] [JJ industrial] [VBG holding] [NN company] [NNP Instituto] [NNP Nacional] [FW de] [NNP Industria]

(9.20) [NP Shearson's] [JJ easy-to-film], [JJ black-and-white] “[SBAR Where We Stand]” [NNS commercials]

Probabilistic Grammar Assumptions

- We're assuming that there is a **grammar** to be used to parse with.
- We're assuming the existence of a large robust **dictionary** with parts of speech
- We're assuming the ability to parse (i.e. a **parser**)
- Given all that... we can parse probabilistically

Typical Approach

- Bottom-up (CKY) dynamic programming approach
- Assign probabilities to constituents as they are completed and placed in the table
- Use the max probability for each constituent going up

What's that last bullet mean?

- Say we're talking about a final part of a parse
 - $S \rightarrow_0 NP_i VP_j$

The probability of the S is...

$$P(S \rightarrow NP VP) * P(NP) * P(VP)$$

The green stuff is already known. We're doing bottom-up parsing

Max

- I said the $P(NP)$ is known.
- What if there are multiple NPs for the span of text in question (0 to i)?
- Take the max (where?)

Problems with PCFGs

- The probability model we're using is just based on the rules in the derivation...
 - Doesn't use the words in any real way
 - Doesn't take into account *where* in the derivation a rule is used

Solution

- Add lexical dependencies to the scheme...
 - Infiltrate the predilections of particular words into the probabilities in the derivation
 - I.e. Condition the rule probabilities on the actual words

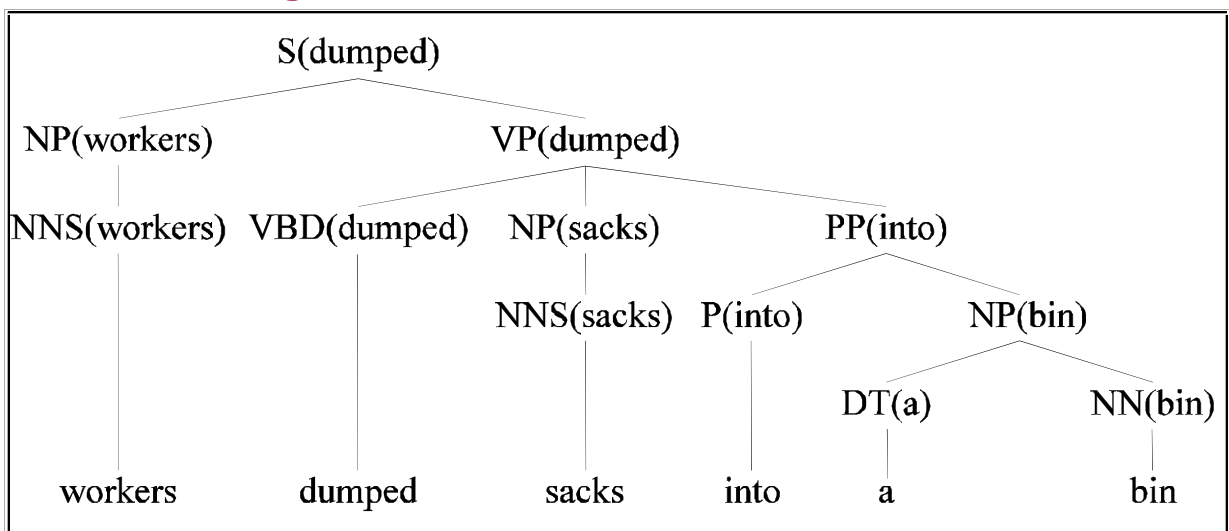
Heads

- To do that we're going to make use of the notion of the head of a phrase
 - The head of an NP is its noun
 - The head of a VP is its verb
 - The head of a PP is its preposition

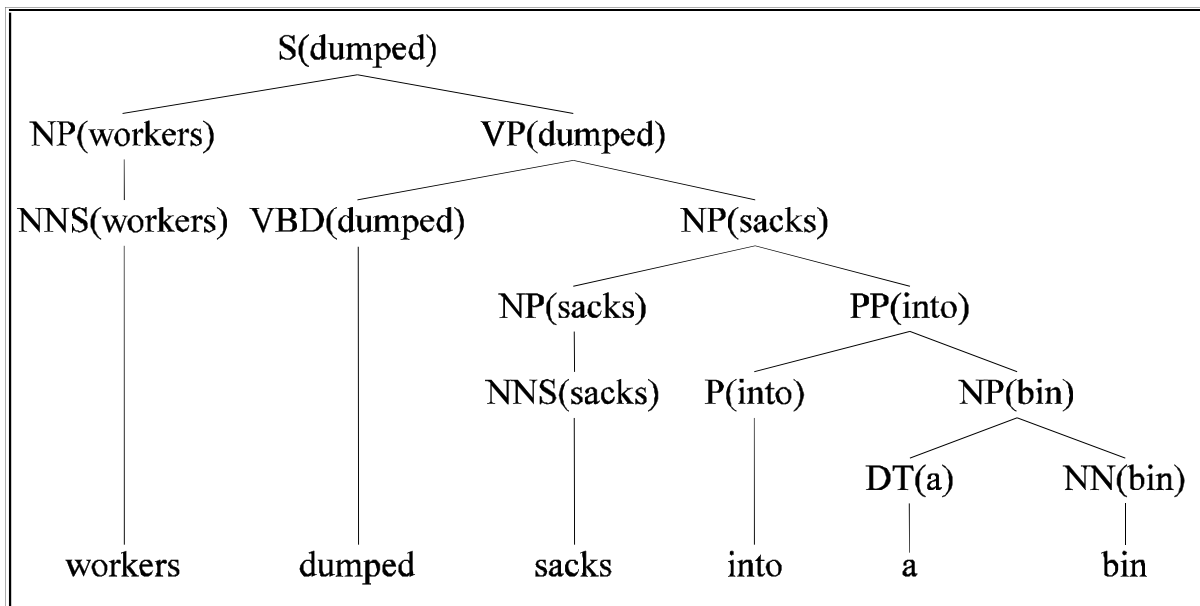
(It's really more complicated than that but this will do.)

Example (right)

Attribute grammar



Example (wrong)



How?

- We used to have
 - $VP \rightarrow V NP PP$ $P(\text{rule}|VP)$
 - That's the count of this rule divided by the number of VPs in a treebank
- Now we have
 - $VP(\text{dumped}) \rightarrow V(\text{dumped}) NP(\text{sacks}) PP(\text{in})$
 - $P(r|VP \wedge \text{dumped is the verb} \wedge \text{sacks is the head of the NP} \wedge \text{in is the head of the PP})$
 - Not likely to have significant counts in any treebank

Declare Independence

- When stuck, exploit independence and collect the statistics you can...
- We'll focus on capturing two things
 - Verb subcategorization
 - Particular verbs have affinities for particular VPs
 - Objects affinities for their predicates (mostly their mothers and grandmothers)
 - Some objects fit better with some predicates than others

Subcategorization

- Condition particular VP rules on their head... so

$r: VP \rightarrow V NP PP \quad P(r|VP)$

Becomes

$P(r | VP \wedge \text{dumped})$

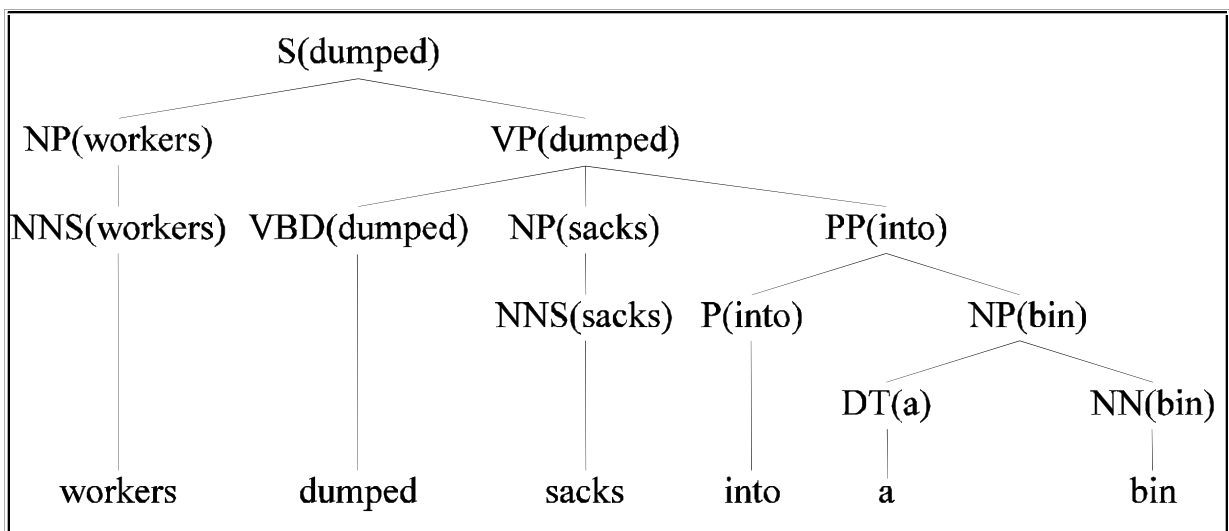
What's the count?

How many times was this rule used with (head) **dump**,
divided by the number of VPs that **dump** appears (as
head) in total

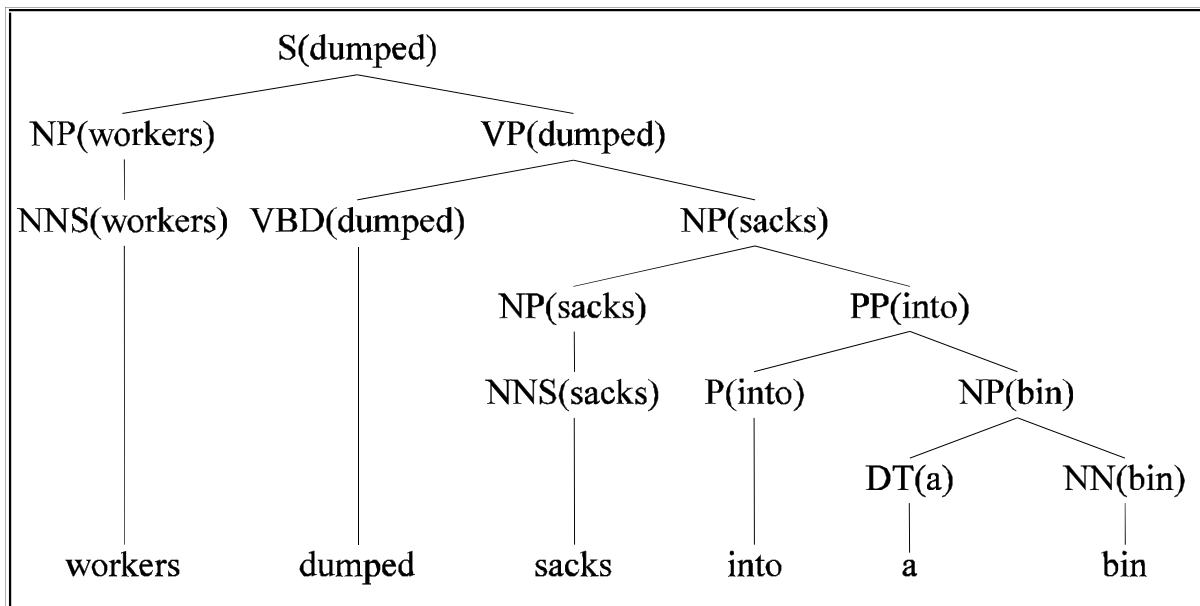
Preferences

- Subcat captures the affinity between VP heads (verbs) and the VP rules they go with.
- What about the affinity between VP heads and the heads of the other daughters of the VP
- Back to our examples...

Example (right)



Example (wrong)



Preferences

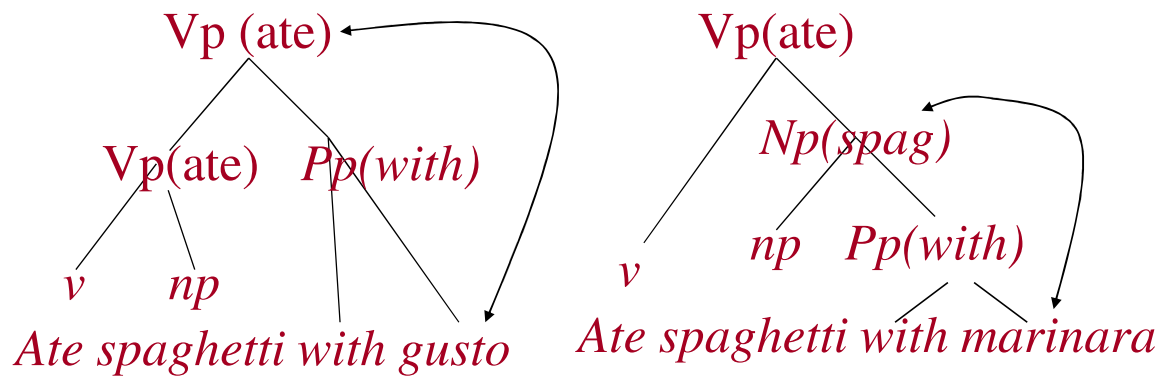
- The issue here is the **attachment** of the PP. So the affinities we care about are the ones between **dumped** and **into** vs. **sacks** and **into**.
- So count the places where **dumped** is the head of a constituent that has a PP daughter with **into** as its head and normalize
- Vs. the situation where **sacks** is a constituent with **into** as the head of a PP daughter.

Preferences (2)

- Consider the VPs
 - Ate spaghetti with gusto
 - Ate spaghetti with marinara
- The affinity of **gusto** for **eat** is much larger than its affinity for **spaghetti**
- On the other hand, the affinity of **marinara** for **spaghetti** is much higher than its affinity for **ate**

Preferences (2)

- Note the relationship here is more distant and doesn't involve a headword since *gusto* and *marinara* aren't the heads of the PPs.



Summary

- Context-Free Grammars
- Parsing
 - Top Down, Bottom Up Metaphors
 - Dynamic Programming Parsers: CKY. Earley
- Disambiguation:
 - PCFG
 - Probabilistic Augmentations to Parsers
 - Treebanks